Managing Enterprise Cybersecurity MIS 4596

Unit#8

Agenda

- Symmetric cryptography
- Block versus Stream ciphers
- Block ciphers
- Block ciphers mode of operations
- Hashes

Symmetric and asymmetric algorithms

...both are 2-way functions that support encryption & decryption

- Symmetric cryptography
  - Use a copied pair of symmetric (identical) secret keys
  - The sender and the receive use the same key for encryption and decryption functions
- Asymmetric cryptography
  - Also know as "public key cryptography"
  - Use different ("asymmetric") keys for encryption and decryption
  - One is called the "private key" and the other is the "public key"

A strong cipher contains

2 main attributes

1. **Confusion:** usually carried out through <u>substitution</u>

2. **Diffusion:** Usually carried out through <u>transposition</u>

Harris, S. and Maymi, F. (2016) <u>All-In-One CISSP Exam Guide</u>, McGraw Hill Education

Symmetric cryptography

Strengths:

- Much faster (less computationally intensive) than asymmetric systems.
- Hard to break if using a large key size.

Weaknesses:

- Requires a secure mechanism to deliver keys properly.
- Each pair of users needs a unique key, so as the number of individuals increases, so does the number of keys, possibly making key management overwhelming.
- Provides confidentiality but not authenticity or nonrepudiation.

Symmetric encryption uses the same keys.

Two types: Stream and Block Ciphers

- **Stream Ciphers** treat the message a stream of bits and performs mathematical functions on each bit individually

- **Block Ciphers** divide a message into blocks of bits and transforms the blocks one at a time

Encrypt message

Decrypt message

Message
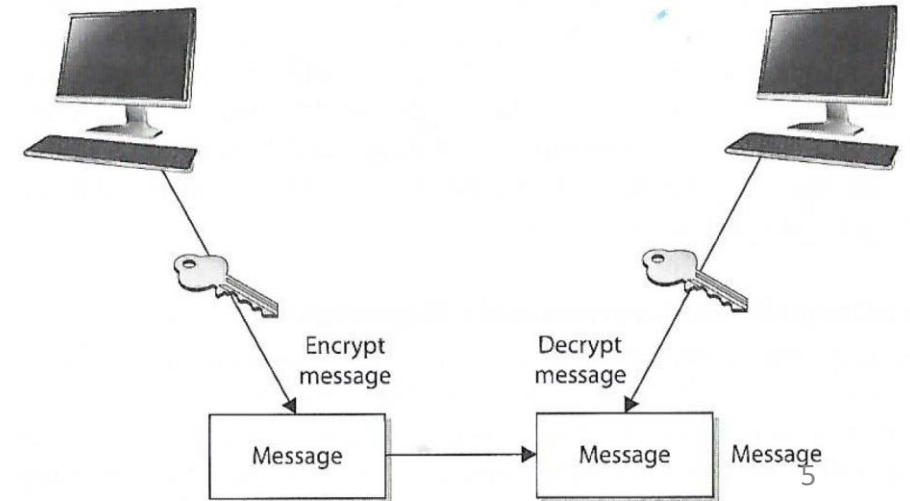
Message

Message

Block Ciphers versus Stream Ciphers

Stream ciphers work on a single bit at a time:

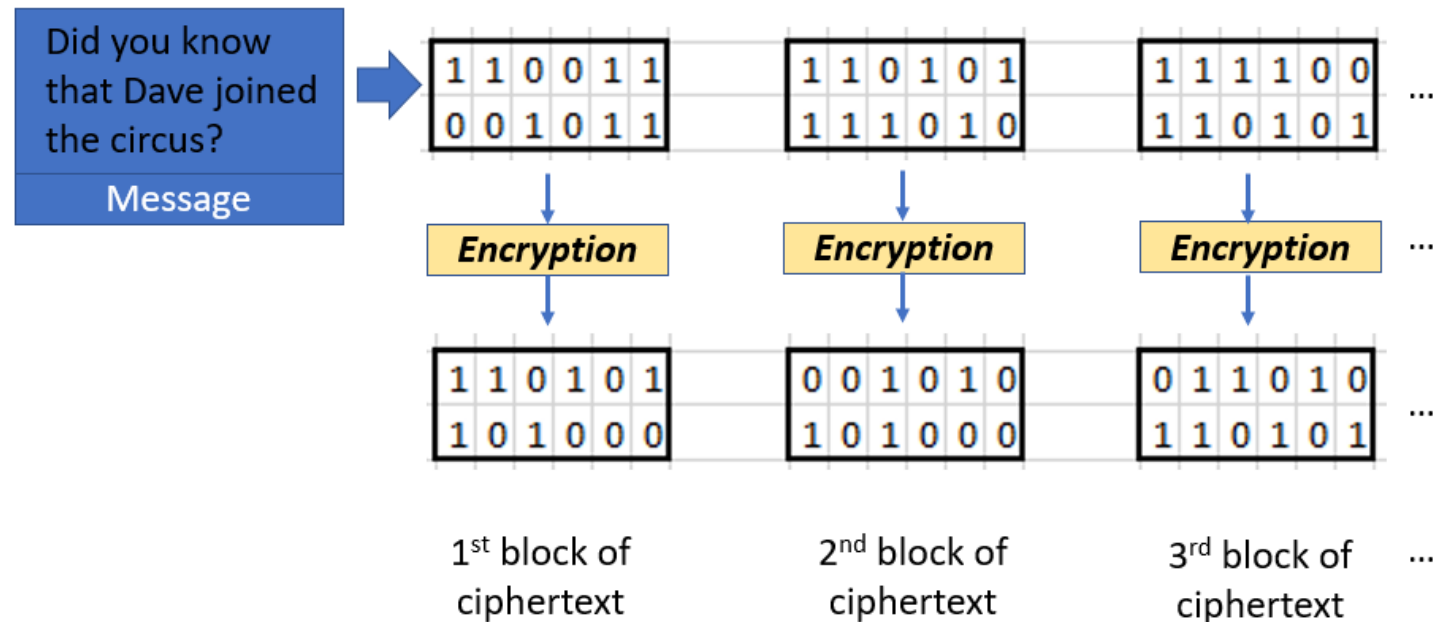| 0	1	1	1	0	0	1	1	Plaintext
			XOR					
1	0	0	1	0	1	0	1	PRNG
1	1	1	0	0	1	1	0	Ciphertext

PRNG = Pseudo Random Number Generator

Block Cyphers ("Cipher")

- Message is divided into blocks of bits
- Blocks are put through encryption functions 1 block at a time

Suppose you are encrypting a 648-bit long message to send to your mother using a block cypher that uses 12 bits

- Your message would be chopped up into 54 blocks each 12 bits long
- Each block, in turn, would be run through a series of encryption functions (substitution and transposition)
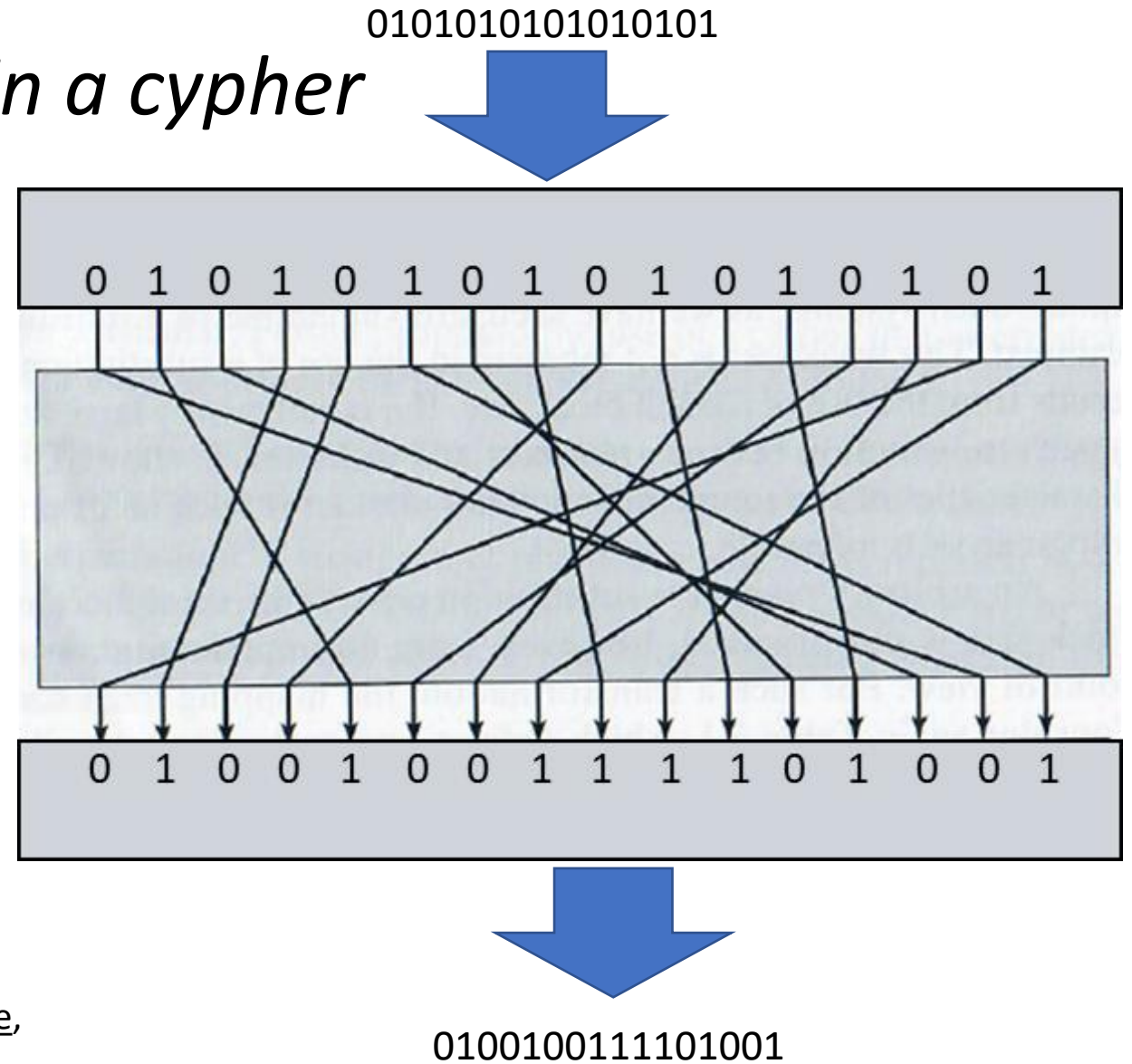- Ending up with 54 blocks of ciphertext

Did you know that Dave joined the circus?
Message

| 1 1 0 0 1 1 | | 1 1 0 1 0 1 | | 1 1 1 1 0 0 | ...
| 0 0 1 0 1 1 | | 1 1 1 0 1 0 | | 1 1 0 1 0 1 |

Encryption   Encryption   Encryption ...

| 1 1 0 1 0 1 | | 0 0 1 0 1 0 | | 0 1 1 0 1 0 | ...
| 1 0 1 0 0 0 | | 1 0 1 0 0 0 | | 1 1 0 1 0 1 |

1st block of ciphertext   2nd block of ciphertext   3rd block of ciphertext ...

Note: Block ciphers typically use 64, 128, 512 bits at a time

Remember this example:

2 main attributes combined in a cypher

0101010101010101

1. Confusion: usually carried out through substitution

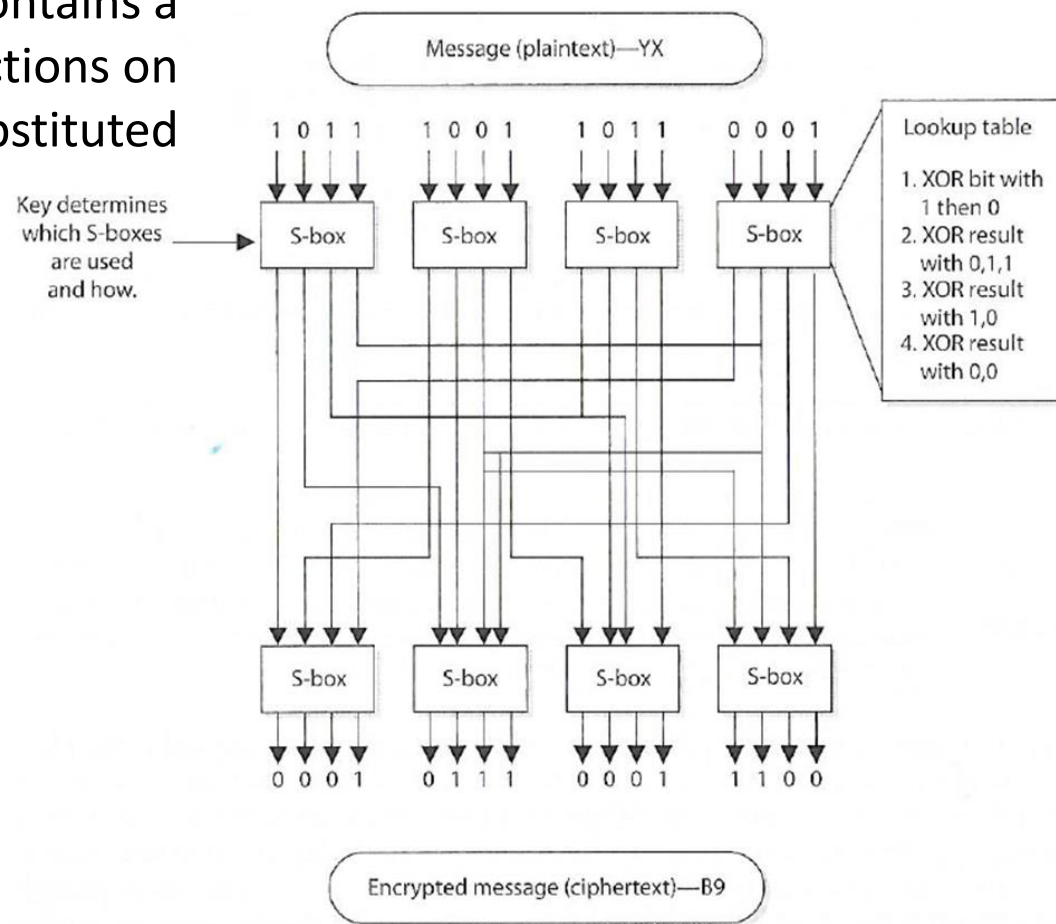2. **Diffusion:** Usually carried out through transposition

| 0 | 1 | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 1 |

| 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 1 | 1 | 0 | 1 | 0 | 0 | 1 |

0100100111101001

Harris, S. and Maymi, F. (2016) All-In-One CISSP Exam Guide, McGraw Hill Education

Block Ciphers versus Stream Ciphers

In contrast, block ciphers encrypt a block of bits at a time

In this example, each Substitution Box (S-box) contains a lookup table used by the algorithm as instructions on how the bits are substituted

Plaintext	Ciphertext
0000	1110
0001	0100
0010	1101
0011	0001
0100	0010
0101	1111
0110	1011
0111	1000
1000	0011
1001	1010
1010	0110
1011	1100
1100	0101
1101	1001
1110	0000
1111	0111

Encryption table

Ciphertext	Plaintext
0000	1110
0001	0011
0010	0100
0011	1000
0100	0001
0101	1100
0110	1010
0111	1111
1000	0111
1001	1101
1010	1001
1011	0110
1100	1011
1101	0010
1110	0000
1111	0101

Decryption table

Message (plaintext)—YX

1 0 1 1   1 0 0 1   1 0 1 1   0 0 0 1

Lookup table

Key determines which S-boxes are used and how.

S-box   S-box   S-box   S-box

1. XOR bit with 1 then 0
2. XOR result with 0,1,1
3. XOR result with 1,0
4. XOR result with 0,0

S-box   S-box   S-box   S-box

0 0 0 1   0 1 1 1   0 0 0 1   1 1 0 0

Encrypted message (ciphertext)—B9

...followed by transposition...

Block Ciphers

- Message is divided into blocks of bits

- Blocks are put through mathematical functions 1 block at a time

You send the message to your mother. She uses the same block cipher and key (symmetric) to decipher the message

- The 54 ciphertext blocks go back through the algorithm in the reverse sequence
- Resulting in your original plaintext message your message

Did you know that Dave joined the circus?

Message

| 1 1 0 0 1 1 |
| 0 0 1 0 1 1 |

| 1 1 0 1 0 1 |
| 1 1 1 0 1 0 |

| 1 1 1 1 0 0 |
| 1 1 0 1 0 1 |

...

Decryption   *Decryption*   *Decryption*   ...

| 1 1 0 1 0 1 |
| 1 0 1 0 0 0 |

| 0 0 1 0 1 0 |
| 1 0 1 0 0 0 |

| 0 1 1 0 1 0 |
| 1 1 0 1 0 1 |

...

1st block of ciphertext   2nd block of ciphertext   3rd block of ciphertext   ...
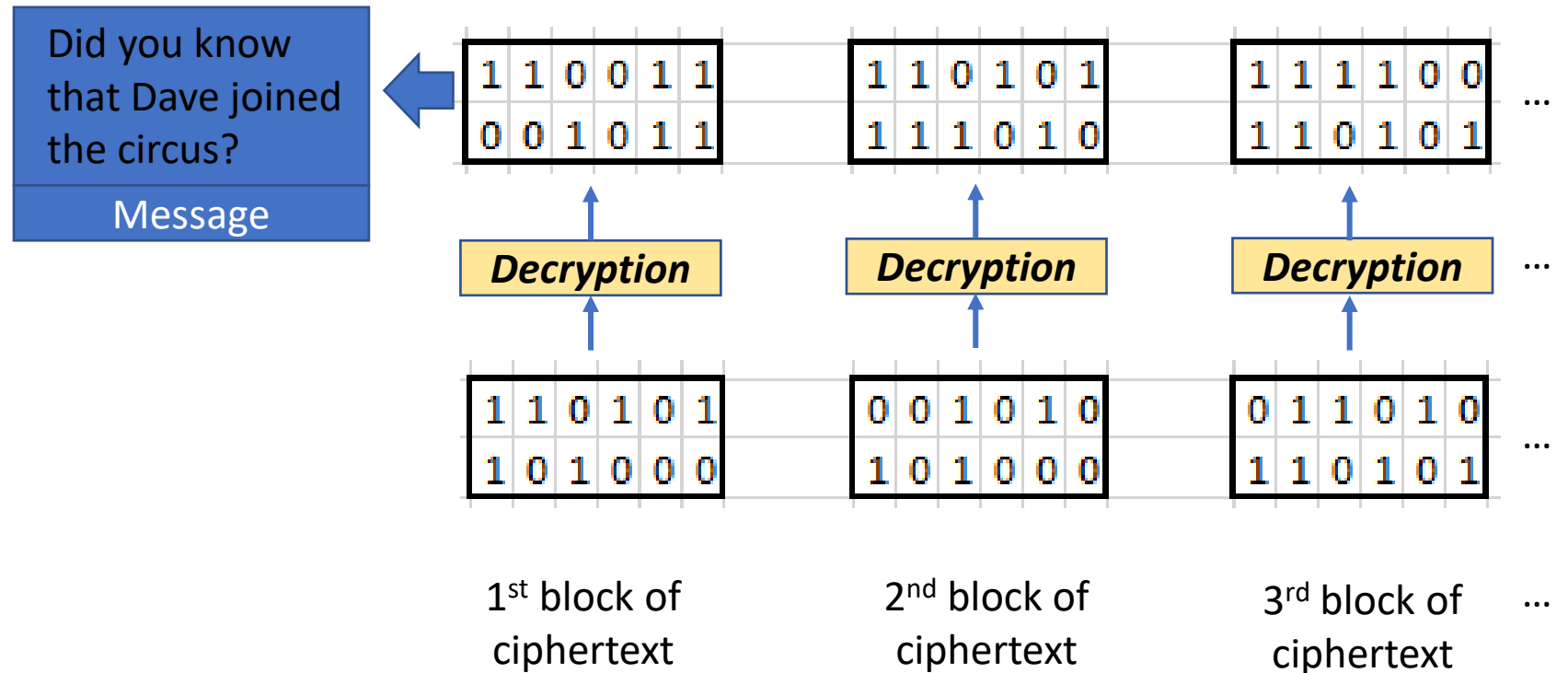
Modern Block Ciphers

- Use block sizes of 128-bits or greater
    - Examples of Block Ciphers that can be used are:
        - AES (NIST's 2001 Advanced Encryption Standard – originally known as Rijndael)
            - 128 bit block size, but 3 different key lengths: 128, 192, and 256 bits
        - Blowfish
        - Twofish
        - Serpent
- Do not use these examples of block ciphers which have a 56 bit key length, which is too small to provide secure encryption:
    - DES (Data Encryption Standard)
    - 3DES

Practical Cryptanalysis

DES Cracker:

- A DES key search machine
- Contains 1,536 chips
- Cost: $250,000
- Searches 88 billion keys per second
- In 1998, Electronic Freedom Foundation (EFF) won $10,000 in RSA Laboratory's "**DES Challenge II-2**" by successfully finding a DES key through brute force in 56 hours
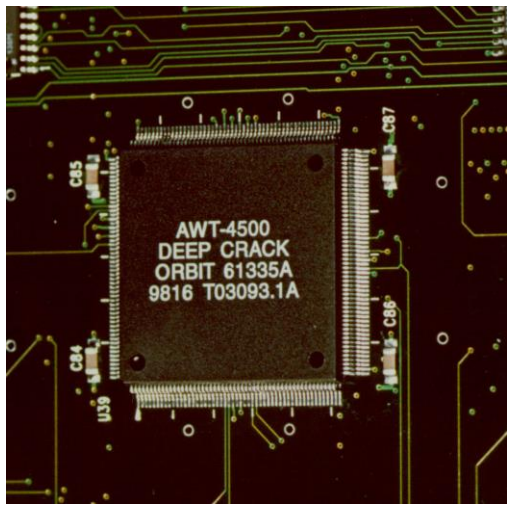
- DES uses a 56-bit key, meaning that there are 256 possible keys under which a message can be encrypted.
  - This is: 72,057,594,037,927,936, or approximately 72 quadrillion possible keys

- One of the major criticisms of DES, when proposed in 1975, was that the key size was too short
  - Martin Hellman and Whitfield Diffie of Stanford University estimated that a machine fast enough to test that many keys in a day would have cost about $20 million in 1976,
    - An affordable sum to national intelligence agencies

- Subsequent advances in the price/performance of chips kept reducing that cost until, twenty years later, it became affordable for even a small nonprofit organization such as the EFF to mount a realistic attack

Agenda

- ✓ Updated Schedule
- ✓ Symmetric cryptography
- ✓ Block versus Stream ciphers
- ✓ Block ciphers
- • Block ciphers mode of operations
- • Hashes

Block cipher's "mode of operation"

5 modes of operation are used to tailor them for use in different applications:

1. **ECB – Electronic Code Book mode**

2. **CBC – Cipher Block Chaining mode**

3. CFB – Cipher FeedBack mode

4. OFB – Output FeedBack mode

5. CTR – CounTeR mode

ECB – Electronic Code Book mode

- A data block of a certain size (e.g. 64 bits or 128 bits or...) is entered into the algorithm with the key, and a block of cipher text is produced

$C_i = Encrypt(Key, P_i)$

for $i = 1, ..., k$

Where:
- *C_i is block i of ciphertext*
- *P_i is a block of plaintext*
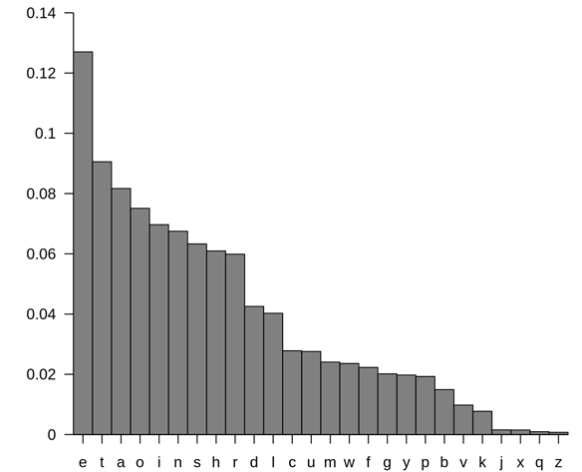
- Encrypts every block <u>the same way every time for a given key</u>

- Why is this a problem?
  - ➢ This is a problem because **frequency analysis** of the encrypted text can reveal a lot of information
  - ➢ Not enough randomness

Solution: CBC – Cipher Block Chaining mode

- Is much more secure

- Does not reveal a pattern of encryption for frequency analysis

- Each block of text, the key, and the value based on the previous block are processed in the algorithm and applied to the next block of text
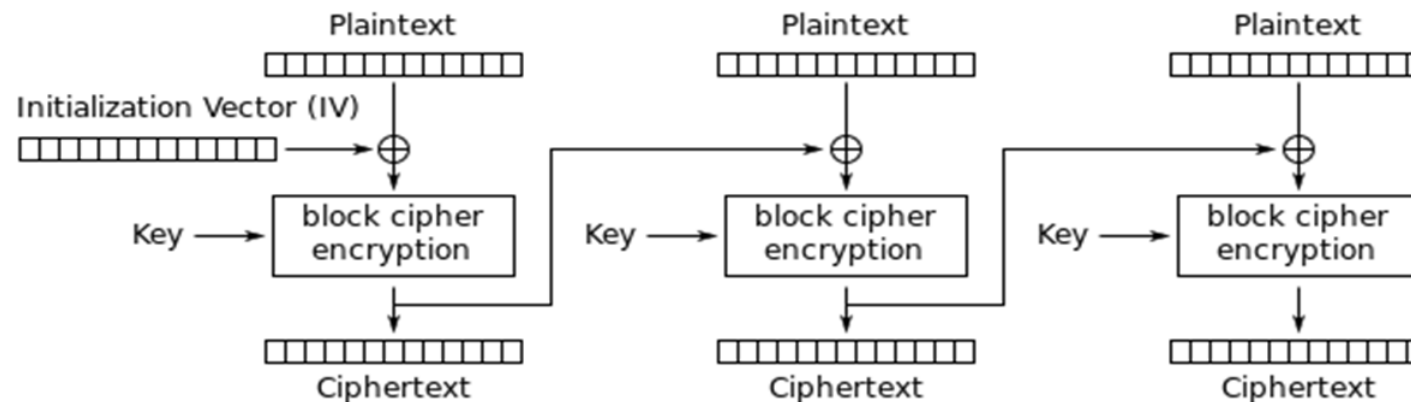
- XORs a plaintext with the **last** encrypted block before encrypting it. This ensures that the same plaintext is encrypted differently every time.

- Requires an initialization vector (or IV) to get started, since the first block doesn't have a previous encrypted block to XOR against.

A similar concept to *diffusion* is known as the

The Avalanche Effect

A changed to a single plaintext bit should have an influence over several of the resulting ciphertext bits

In a strong block cipher, if 1 plaintext bit is changed, it will change every ciphertext bit with the probability of 50%
    That is, if 1 plaintext bit changes, then about ½ of the ciphertext bits will change

     Avalanche Effect: A small change to the key or to the plaintext should cause drastic changes to the resulting ciphertext

Harris, S. and Maymi, F. (2019) <u>All-In-One CISSP Exam Guide</u>, McGraw Hill Education

Visa

Original Image

Block cipher with ECB
(Electronic Code Book)
encryption

Not good!
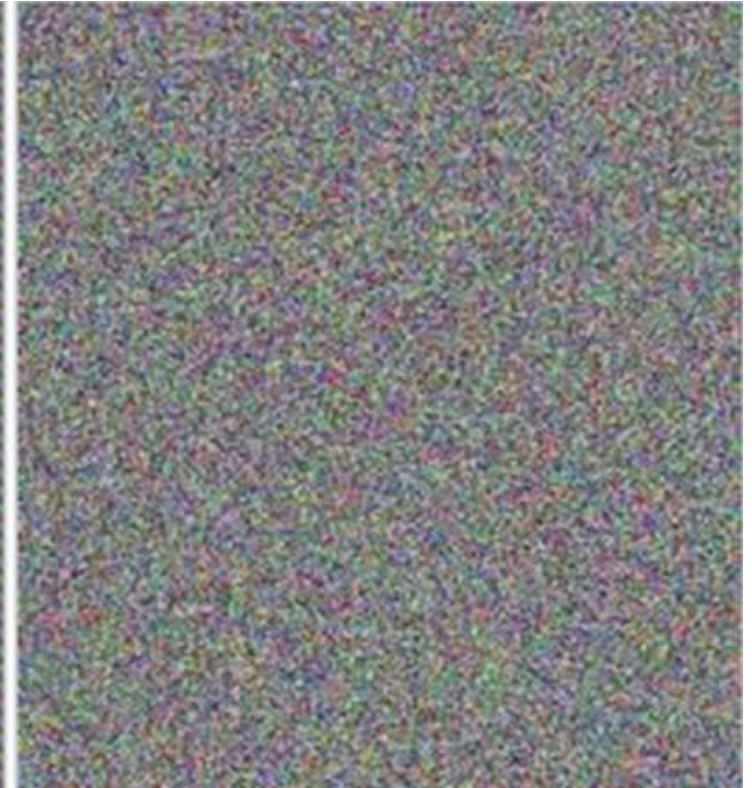
Block cipher with CBC
(Cipher Block Chaining) or any
of the other modes of
encryption

These are good!

Agenda

✓Updated Schedule

✓Symmetric cryptography

✓Block versus Stream ciphers

✓Block ciphers

✓Block ciphers mode of operations

• Hashes

One-way Hash function

- Assures message **integrity**

- A function that takes a variable-length string (i.e. message) and produces a fixed-length value called a hash value

- Does not use keys

1. Sender puts message through hashing function
2. Message digest generated
3. Message digest appended to the message
4. Sender sends message to receiver
5. Receiver puts message through hashing function
6. Receiver generates message digest value
7. Receiver compares the two message digests values. If they are the same, the message has not been altered
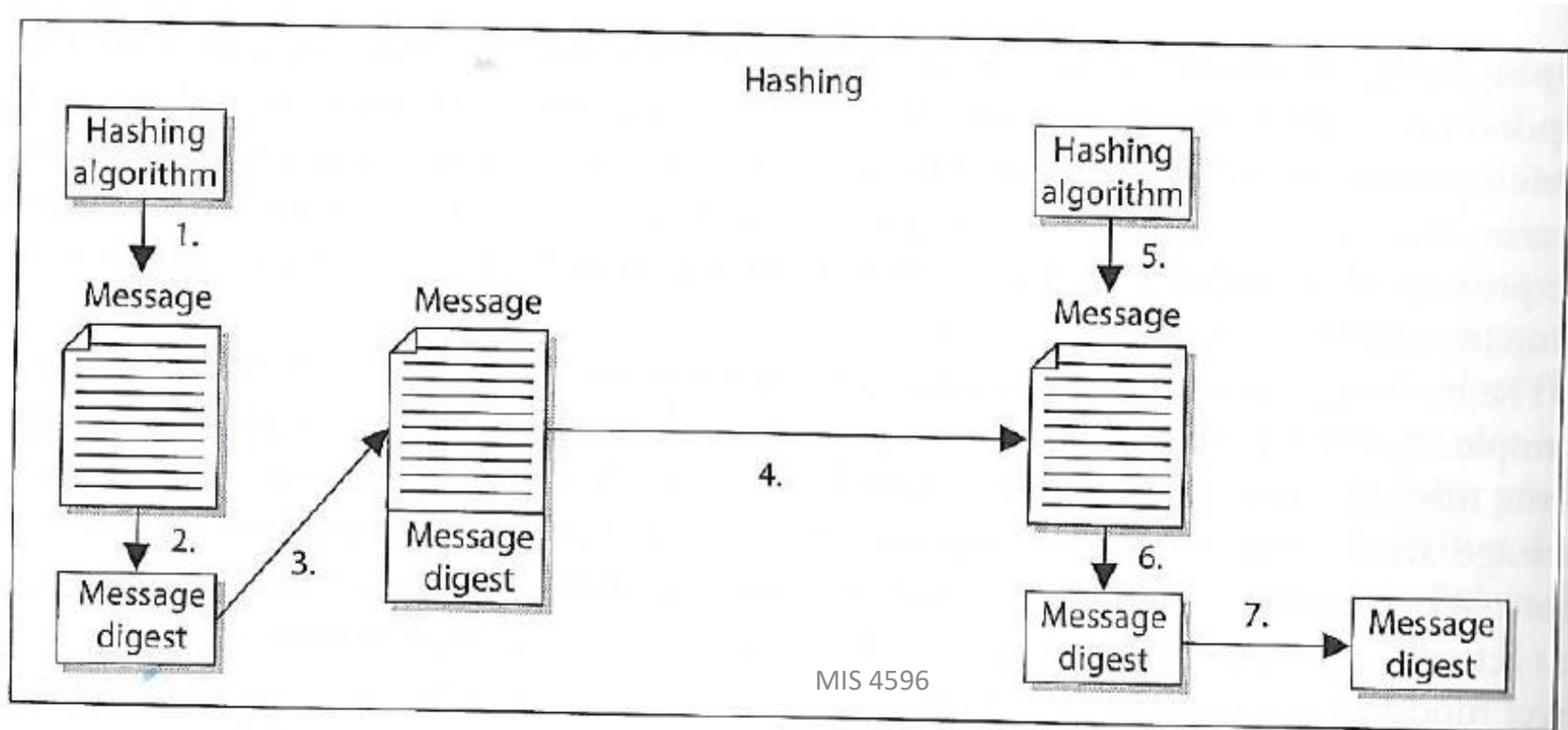
Hashing results in **fixed-sized output**

Names for the output of a hashing functions include "hash" and a *message digest* **(md)**, because a hash "digests" an input of any size down to a **fixed-sized output**

- No matter the size of the input, the output is the same, for example the md5 hash function's output:
  - Letter 'a' in binary: 01000001 => md5 hash => 32-character string
  - Blu-ray disk digest   => md5 hash => 32-character string
  - 6 TB hard drive digest   => md5 hash => 32-character string

One-way hash example…

One useful example: Testing the integrity of a file (e.g. program) downloaded from the internet…

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Image Name	Download	Size	Version	sha256sum
Kali 64 bit	HTTP | Torrent	2.8G	2017.2	4556775bfb981ae64a3cb19aa0b73e8dcac6e4ba524f31c4bc14c9137b99725d
Kali 32 bit	HTTP | Torrent	2.9G	2017.2	7f5000d8f55469264399a8bb7358fc22bec87fb1dc8a51b87f26876634e3effc
Kali 64 bit Light	HTTP | Torrent	0.8G	2017.2	369a29deff40dfff4f53fb47a6015d41d4ada8833a0b6e159657d2f223670f8b
Kali 32 bit Light	HTTP | Torrent	0.8G	2017.2	f6ee21b2880501caee8aa47960e8f424dab5fae1a13ba4b4e02d45152b6acd0d
Kali 64 bit e17	HTTP | Torrent	2.6G	2017.2	20dee81d9891aa6dcfe505a68692f98f981b43a14234d18d9edd92373d6ed6ab
Kali 64 bit Mate	HTTP | Torrent	2.8G	2017.2	9c99a2cc52b1d48875681d12e1fcf6b0b003d44f7ceb610438b5bea148414810
Kali 64 bit Xfce	HTTP | Torrent	2.7G	2017.2	9ecf6a054de1e3ad04d4063e3d347efb31326078c104ec2e78ab456fc4d2a578
Kali 64 bit LXDE	HTTP | Torrent	2.7G	2017.2	c832df6b7a8e7074a5d7f5a50245b840a3df72ffdd4d19a5d1f647beebb4f299
Kali armhf	HTTP | Torrent	0.6G	2017.2	a7f3e648ce9784589245c18d84e2273eb1f4ec1b78244a2c6d4465f3744c9198

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One-way hash example…

Testing the integrity of a file (e.g. program) from the internet…

Image Name	Download	Size	Version	sha256sum
Kali 64 bit	HTTP | Torrent	2.8G	2017.2	4556775bfb981ae64a3cb19aa0b73e8dcac6e4ba524f31c4bc14c9137b99725d
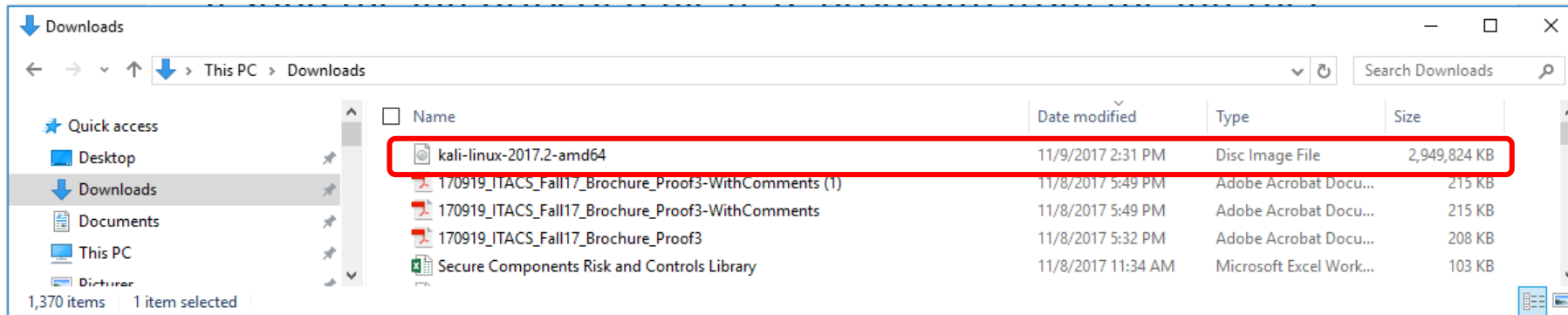
Is the Kali I downloaded the same Kali that was published?

One-way hash example…

One-way hash example...

Example 1: Compute the hash value for a PowerShell.exe file

```
PS C:\> Get-FileHash $pshome\powershell.exe | Format-List

Algorithm : SHA256
Hash      : 6A785ADC0263238DAB3EB37F4C185C8FBA7FEB5D425D034CA9864F1BE1C1B473
Path      : C:\Windows\System32\WindowsPowerShell\v1.0\powershell.exe
```

This command uses the **Get-FileHash** cmdlet to compute the hash value for the Powershell.exe file. The hash algorithm used is the default, SHA256. The output is piped to the Format-List cmdlet to format the output as a list.

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One-way hash example…

Image Name	Download	Size	Version	sha256sum
Kali 64 bit	HTTP | Torrent	2.8G	2017.2	4556775bfb981ae64a3cb19aa0b73e8dcac6e4ba524f31c4bc14c9137b99725d


```
Windows PowerShell                                                    —  □  ×

PS C:\Users\tue87168> cd Downloads
PS C:\Users\tue87168\Downloads> dir *.iso


    Directory: C:\Users\tue87168\Downloads


Mode                LastWriteTime         Length Name
----                -------------         ------ ----
-a----        8/10/2017   10:55 AM      674803712 CSET_8.0 (1).iso
-a----        8/10/2017   11:03 AM      674803712 CSET_8.0 (2).iso
-a----        6/12/2017   10:29 AM      674803712 CSET_8.0.iso
-a----        9/27/2017    3:03 PM     2421987328 en_project_professional_2016_x86_x64_dvd_6962236.iso
-a----        10/3/2017    8:49 PM     2421987328 en_visio_professional_2016_x86_x64_dvd_6962139.iso
-a----       11/11/2016   11:45 AM     1469054976 Fedora-Live-Workstation-x86_64-23-10.iso
-a----        11/9/2017    2:31 PM     3020619776 kali-linux-2017.2-amd64.iso


PS C:\Users\tue87168\Downloads> Get-FileHash kali-linux-2017.2-amd64.iso | Format-List


Algorithm : SHA256
Hash      : 4556775BFB981AE64A3CB19AA0B73E8DCAC6E4BA524F31C4BC14C9137B99725D
Path      : C:\Users\tue87168\Downloads\kali-linux-2017.2-amd64.iso


PS C:\Users\tue87168\Downloads> _
```
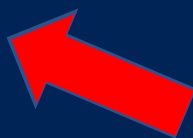
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One-way hash example…

Windows PowerShell

```
PS C:\Users\tue87168\Downloads> dir *.txt

    Directory: C:\Users\tue87168\Downloads


Mode                LastWriteTime         Length Name
----                -------------         ------ ----
-a----         11/9/2017    3:04 PM            15 MIS5206-IsGood.txt


PS C:\Users\tue87168\Downloads> type MIS5206-IsGood.txt
MIS5206 is good
PS C:\Users\tue87168\Downloads> Get-FileHash MIS5206-IsGood.txt | Format-List


Algorithm : SHA256
Hash      : E6F053ADE3857C0EDC2896B229D0B91D4752B2D9D8C9BD4B2A45A4ACCB3999DD
Path      : C:\Users\tue87168\Downloads\MIS5206-IsGood.txt



PS C:\Users\tue87168\Downloads> type MIS5206-IsGood.txt
MIS5206 is goop
PS C:\Users\tue87168\Downloads> Get-FileHash MIS5206-IsGood.txt | Format-List


Algorithm : SHA256
Hash      : 877B45EA5D40D98FF8D1ABD919E154F446FEA11387DBB13DDEE448F9932928A5
Path      : C:\Users\tue87168\Downloads\MIS5206-IsGood.txt



PS C:\Users\tue87168\Downloads>
```

Notice the amount of confusion and diffusion resulting from a 1-character change!

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File Integrity Monitoring

CHANGE LOGS

EMAILS

SERVER STACK

FIM MANAGEMENT
SERVER
[Records & scans
important files
for modification]

REAL-TIME
NOTIFICATION
[of Suspicious
Behavior]

- An internal control process

- Validates the integrity of operating system and application software files

- Uses hash verification to compare the current file state and a known good baseline state

- Involves calculating and storing a hash value of a known good version of the file ("original baseline")

- Compares the baseline with the calculated hash of the current state of the file to detect unauthorized changes

How are passwords stored in Linux?

```
cat /etc/passwd
root:x:0:0:root:/root:/bin/bash
daemon:x:1:1:daemon:/usr/sbin:/bin/sh
bin:x:2:2:bin:/bin:/bin/sh
sys:x:3:3:sys:/dev:/bin/sh
sync:x:4:65534:sync:/bin:/bin/sync
games:x:5:60:games:/usr/games:/bin/sh
man:x:6:12:man:/var/cache/man:/bin/sh
lp:x:7:7:lp:/var/spool/lpd:/bin/sh
mail:x:8:8:mail:/var/mail:/bin/sh
news:x:9:9:news:/var/spool/news:/bin/sh
uucp:x:10:10:uucp:/var/spool/uucp:/bin/sh
proxy:x:13:13:proxy:/bin:/bin/sh
www-data:x:33:33:www-data:/var/www:/bin/sh
backup:x:34:34:backup:/var/backups:/bin/sh
list:x:38:38:Mailing List Manager:/var/list:/bin/sh
irc:x:39:39:ircd:/var/run/ircd:/bin/sh
gnats:x:41:41:Gnats Bug-Reporting System (admin):/var/lib/gnats:/bin/sh
nobody:x:65534:65534:nobody:/nonexistent:/bin/sh
libuuid:x:100:101::/var/lib/libuuid:/bin/sh
dhcp:x:101:102::/nonexistent:/bin/false
syslog:x:102:103::/home/syslog:/bin/false
klog:x:103:104::/home/klog:/bin/false
sshd:x:104:65534::/var/run/sshd:/usr/sbin/nologin
msfadmin:x:1000:1000:msfadmin,,,:/home/msfadmin:/bin/bash
bind:x:105:113::/var/cache/bind:/bin/false
postfix:x:106:115::/var/spool/postfix:/bin/false
ftp:x:107:65534::/home/ftp:/bin/false
postgres:x:108:117:PostgreSQL administrator,,,:/var/lib/postgresql:/bin/bash
mysql:x:109:118:MySQL Server,,,:/var/lib/mysql:/bin/false
tomcat55:x:110:65534::/usr/share/tomcat5.5:/bin/false
distccd:x:111:65534::/:/bin/false
user:x:1001:1001:just a user,111,,:/home/user:/bin/bash
service:x:1002:1002:,,,:/home/service:/bin/bash
telnetd:x:112:120::/nonexistent:/bin/false
proftpd:x:113:65534::/var/run/proftpd:/bin/false
statd:x:114:65534::/var/lib/nfs:/bin/false
```

```
cat /etc/shadow
root:$1$/avpfBJ1$x0z8w5UF9Iv./DR9E9Lid.:14747:0:99999:7:::
daemon:*:14684:0:99999:7:::
bin:*:14684:0:99999:7:::
sys:$1$fUX6BPOt$Miyc3UpOzQJqz4s5wFD9l0:14742:0:99999:7:::
sync:*:14684:0:99999:7:::
games:*:14684:0:99999:7:::
man:*:14684:0:99999:7:::
lp:*:14684:0:99999:7:::
mail:*:14684:0:99999:7:::
news:*:14684:0:99999:7:::
uucp:*:14684:0:99999:7:::
proxy:*:14684:0:99999:7:::
www-data:*:14684:0:99999:7:::
backup:*:14684:0:99999:7:::
list:*:14684:0:99999:7:::
irc:*:14684:0:99999:7:::
gnats:*:14684:0:99999:7:::
nobody:*:14684:0:99999:7:::
libuuid:!:14684:0:99999:7:::
dhcp:*:14684:0:99999:7:::
syslog:*:14684:0:99999:7:::
klog:$1$f2ZVMS4K$R9XkI.CmLdHhdUE3X9jqP0:14742:0:99999:7:::
sshd:*:14684:0:99999:7:::
msfadmin:$1$XN10Zj2c$Rt/zzCW3mLtUWA.ihZjA5/:14684:0:99999:7:::
bind:*:14685:0:99999:7:::
postfix:*:14685:0:99999:7:::
ftp:*:14685:0:99999:7:::
postgres:$1$Rw35ik.x$MgQgZUuO5pAoUvfJhfcYe/:14685:0:99999:7:::
mysql:!:14685:0:99999:7:::
tomcat55:*:14691:0:99999:7:::
distccd:*:14698:0:99999:7:::
user:$1$HESu9xrH$k.o3G93DGoXIiQKkPmUgZ0:14699:0:99999:7:::
service:$1$kR3ue7JZ$7GxELDupr5Ohp6cjZ3Bu//:14715:0:99999:7:::
telnetd:*:14715:0:99999:7:::
proftpd:!:14727:0:99999:7:::
statd:*:15474:0:99999:7:::
```

Shadow file

Password file

Cryptanalysis Attacks

- Brute force
  - Trying all key values in the keyspace
- Frequency Analysis
  - Guess values based on frequency of occurrence
- Dictionary Attack
  - Find plaintext based on common words
- Known Plaintext
  - Format or content of plaintext available
- Chosen Plaintext
  - Attack can encrypt chosen plaintext
- Chosen Ciphertext
  - Decrypt known ciphertext to discover key

- Random Number Generator (RNG) Attack
  - Predict initialization vector used by an algorithm
- Social Engineering
  - Humans are the weakest link

MIS 4596

Cryptanalysis Attacks

Collisions

- *Two different messages with the same hash value*
- *Hash algorithms should be resistant to this attack*

```
PS C:\Users\David Lanter\Downloads\MD5-Hash-Collision-Example> get-filehash ProgramA.exe -Algorithm MD5

Algorithm       Hash                                                            Path
---------       ----                                                            ----
MD5             CDC47D670159EEF60916CA03A9D4A007                                C:\Users\David Lanter\Downloads\M...


PS C:\Users\David Lanter\Downloads\MD5-Hash-Collision-Example> get-filehash ProgramB.exe -Algorithm MD5

Algorithm       Hash                                                            Path
---------       ----                                                            ----
MD5             CDC47D670159EEF60916CA03A9D4A007                                C:\Users\David Lanter\Downloads\M...


PS C:\Users\David Lanter\Downloads\MD5-Hash-Collision-Example>
```

```
PS C:\Users\David Lanter\Downloads\MD5-Hash-Collision-Example> Get-Filehash ProgramA.exe -Algorithm SHA512 |Format-List


Algorithm : SHA512
Hash      : 7D6BFC7DE9D16842E48A138EA9FC395BD1994085DAB055225D14E6A7CAE5200BB8FB1C407F5EE35E83939864C762E3154680709C04194F3
            3E37951298C795881
Path      : C:\Users\David Lanter\Downloads\MD5-Hash-Collision-Example\ProgramA.exe



PS C:\Users\David Lanter\Downloads\MD5-Hash-Collision-Example> Get-Filehash ProgramB.exe -Algorithm SHA512 |Format-List


Algorithm : SHA512
Hash      : 383BB26E68F8B7A362CCEB03F66E586F9EEDAA49D900449A0E33BF0359892D4B978AA65F587D69BB0089F91D848DF91774AC973AC22D2D7
            C1ECBD73A7CE77CA2
Path      : C:\Users\David Lanter\Downloads\MD5-Hash-Collision-Example\ProgramB.exe
```

Cryptanalysis Attacks

Collisions

- *Two different messages with the same hash value*
- Based on the "birthday paradox"
- Hash algorithms should be resistant to this attack

The birthday paradox, also known as the birthday problem, states that in a random group of 23 people, there is about a 50 percent chance that two people have the same birthday.

Is the Birthday Attack Real?

There are multiple reasons why this seems like a paradox
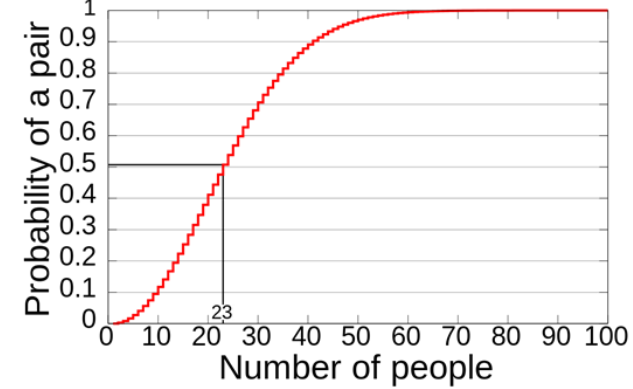
One is that when in a room with 22 other people, if a person compares her/his birthday with the birthdays of the other people it would make for only 22 comparisons—only 22 chances for people to share the same birthday.

When all 23 birthdays are compared against each other, it makes for much more than 22 comparisons. How much more?

Well, the first person has 22 comparisons to make, but the second person was already compared to the first person, so there are only 21 comparisons to make.

The third person then has 20 comparisons, the fourth person has 19 and so on.

If you add up all possible comparisons (22 + 21 + 20 + 19 + … +1) the sum is 253 comparisons, or combinations. Consequently, each group of 23 people involves 253 comparisons, or 253 chances for matching birthdays.

MIS 4596

MD5 (Message Digest 5)

- A 128-bit hash algorithm, still in common use
- Has been broken
- 128-bit hash, but only need $2^{128/2} = 2^{64}$ to find a collision
- Not strong enough for modern computers

SHA -1 (Security Hash Algorithm 1)

- A 160-bit hash algorithm, still in common use
- Has been broken
- 160-bit hash, but only need $2^{160/2} = 2^{80}$ to find a collision
- No longer strong enough for modern computers

Windows PowerShell

PS C:\Users\Dave\Desktop\MD5-Hash-Collision-Example> get-filehash ProgramA.exe -Algorithm SHA256

Algorithm	Hash	Path
SHA256	60D13913155644883F130B85EB24D778314014C9479AEDB5F6323BF38AD3A451	C:\Users\Dave\Desktop\MD5-Hash-Collision-Example\ProgramA.exe

PS C:\Users\Dave\Desktop\MD5-Hash-Collision-Example> get-filehash ProgramB.exe -Algorithm SHA256

Algorithm	Hash	Path
SHA256	1316543942A8C6CD754855500CD37068EDBBD8B31C4979D2825A4E799FED6102	C:\Users\Dave\Desktop\MD5-Hash-Collision-Example\ProgramB.exe

```
Hello, world!

(press enter to quit)
```
ProgramA run

```
This program is evil!!!
Erasing hard drive...1Gb...2Gb... just kidding!
Nothing was erased.

(press enter to quit)_
```
ProgramB run

Windows PowerShell

PS C:\Users\Dave\Desktop\MD5-Hash-Collision-Example> get-filehash ProgramA.exe -Algorithm MD5

Algorithm	Hash	Path
MD5	CDC47D670159EEF60916CA03A9D4A007	C:\Users\Dave\Desktop\MD5-Hash-Collision-Example\ProgramA.exe

PS C:\Users\Dave\Desktop\MD5-Hash-Collision-Example> get-filehash ProgramB.exe -Algorithm MD5

Algorithm	Hash	Path
MD5	CDC47D670159EEF60916CA03A9D4A007	C:\Users\Dave\Desktop\MD5-Hash-Collision-Example\ProgramB.exe

The malware Flame used a MD5 hash collision to hijack Microsoft Windows Update and spread itself across networks

- Flame collected **audio, keystrokes, screenshots** which it sent to a malicious server
- Found a collision within a single millisecond
- Cost ~$200k computing time just for 1ms
- Attributed to advanced persistent threat group Equation Group
- Espionage attacks on countries in and around Iran
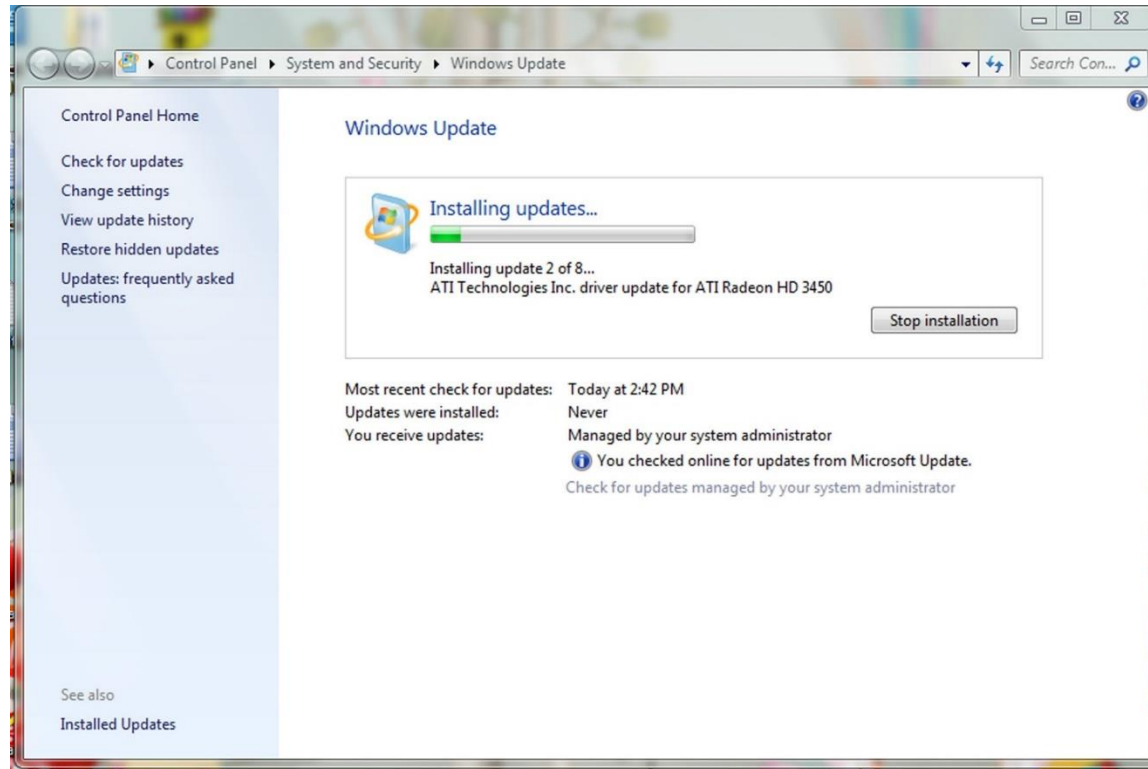
PWN (verb)

1. An act of dominating an opponent.

2. Great, ingenious; applied to methods and objects.

- Originally dates back to the days of WarCraft, when a map designer mispelled "Own" as "Pwn"

- What was originally suppose to be "player has been owned." was "player has been pwned"

Use of the term "Pwn' grew and is now used throughout the online world, especially in online games:
   1. "I pwn these guys on battlenet"

   2. "This strategy pwns!" or "This game pwn."

   https://www.urbandictionary.com/define.php?term=pwn

Equation Group's Flame malware won 2012 "Epic 0wnage" Pwnie award

- Pwnie Awards recognize both excellence and incompetence in the field of information security
- Awards are presented yearly at the Black Hat Security Conference

Hashing algorithms are used for browser ssl (secure sockets layer)

- In 2014, many sites were still using SHA-1, at the time known to be dangerously vulnerable

- Google declared state of emergency to push companies to upgrade

Why Google is Hurrying the Web to Kill SHA-1

published by Eric Mill on September 7, 2014, 58 comments
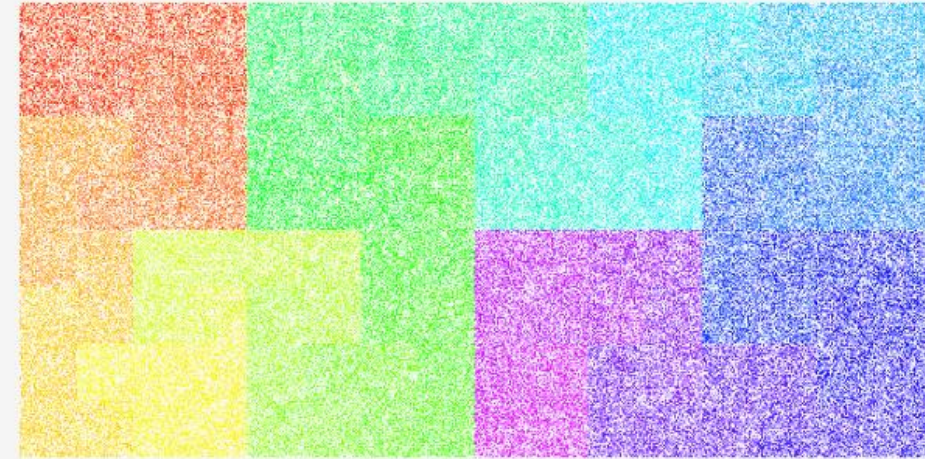
Hilbert map of hashing algorithms, by Ian Boyd

Most of the secure web is using an insecure algorithm, and Google's just declared it to be a slow-motion emergency.

Something like 90% of websites that use SSL encryption — 🔒 https:// — use an algorithm called SHA-1 to protect themselves from being impersonated. This guarantees that when you go to 🔒 https://www.facebook.com , you're visiting the real Facebook and not giving your password to an attacker.

Unfortunately, **SHA-1 is dangerously weak**, and has been for a long time. It gets weaker every year, but remains widely used on the internet. Its replacement, SHA-2, is strong and supported just about everywhere.

Google recently announced that if you use Chrome, then you're about to start seeing a progression of warnings for many secure websites:

🔒 https://   ➡   📄 https://   ➡   ✖ **https**://
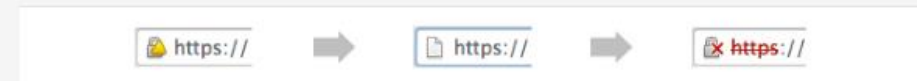
What's about to befall websites with SHA-1 certificates that expire in 2017, in Chrome.

The first set of warnings will hit before Christmas, and will keep getting more stern over the

SHA-2 uses 224, **256**, 384, and 512-bit hashes
- But… it is built using the design of SHA-1, and prone to the same weaknesses
- It's believed to be a matter of time before SHA-2 is also exploited

- SHA-3 was just ratified recently by NIST, the U.S. National Institute of Standards and Technology
  - It was the result of a six-year hashing competition. Also uses 224-, 256-, 384-, 512-bit hashes

Why does this matter for businesses?

Business needs a reliable way to prove integrity of data, files, programs, that can be trusted

Agenda

- ✓ Updated Schedule
- ✓ Symmetric cryptography
- ✓ Block versus Stream ciphers
- ✓ Block ciphers
- ✓ Block ciphers mode of operations
- ✓ Hashes