# Managing Enterprise Cybersecurity MIS 4596

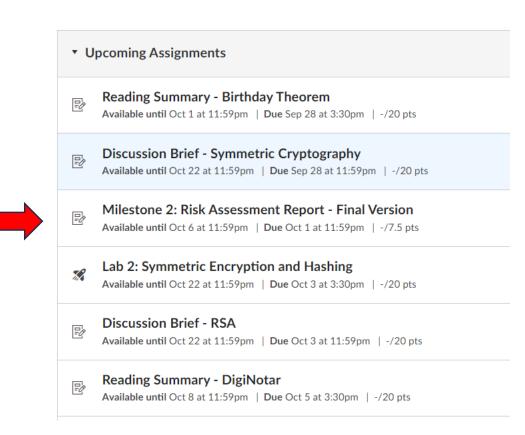
Unit#8

### Agenda

- Updated Schedule
- Symmetric cryptography
- Block versus Stream ciphers
- Block ciphers
- Block ciphers mode of operations
- Hashes

#### Updated Schedule

DATES	TOPICS				
Tuesday, Aug 29	Introduction to the Course				
Thursday, Aug 31	Threat modeling				
Tuesday, Sep 5	Risk Assessment				
Thursday, Sep 7	7 Class Cancelled				
Tuesday, Sep 12	Milestone 1 Risk Assessment Report – Q&A, Introduction to Google Cloud Platform (GCP) & Linux				
Thursday, Sep 14	Data Privacy				
Sunday, Sep 17	Milestone 1 Risk Assessment Report Due				
Tuesday, Sep 19	Introduction to Cryptography				
Thursday, Sep 21	Introduction to Cryptographycontinued				
Tuesday, Sep 26	Milestone 1 Report Feedback				
Thursday, Sep 28	Symmetric Cryptography & Hashing				
Sunday, Oct 1	Milestone 2 Final Risk Assessment Report Due				
Tuesday, Oct 3	Asymmetric Cryptography & Digital Signatures				
Thursday, Oct 5	Digital Certificates and Public Key Infrastructures				



#### Symmetric and asymmetric algorithms

#### ...both are 2-way functions that support encryption & decryption

- Symmetric cryptography
  - Use a copied pair of symmetric (identical) secret keys
  - The sender and the receive use the same key for encryption and decryption functions
- Asymmetric cryptography
  - Also know as "public key cryptography"
  - Use different ("asymmetric") keys for encryption and decryption
  - One is called the "private key" and the other is the "public key"

### A strong cipher contains

2 main attributes

- 1. Confusion: usually carried out through <u>substitution</u>
- 2. Diffusion: Usually carried out through transposition

# Symmetric cryptography

#### Strengths:

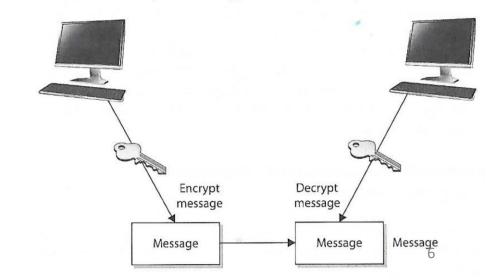
- Much faster (less computationally intensive) than asymmetric systems.
- Hard to break if using a large key size.

#### Weaknesses:

- Requires a secure mechanism to deliver keys properly.
- Each pair of users needs a unique key, so as the number of individuals increases, so does the number of keys, possibly making key management overwhelming.
- Provides confidentiality but not authenticity or nonrepudiation.

#### Two types: Stream and Block Ciphers

- Stream Ciphers treat the message a stream of bits and performs mathematical functions on each bit individually
- Block Ciphers divide a message into blocks of bits and transforms the blocks one at a time MIS 4596



Symmetric encryption uses the same keys.

### Block Ciphers versus Stream Ciphers

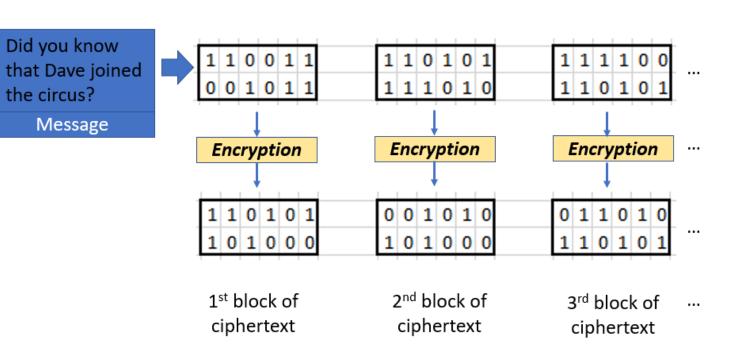
Stream ciphers work on a single bit at a time:

									11
0	1	1	1	0	0	1	1	Plaintext	
			XC	DR					I
1	0	0	1	0	1	0	1	PRNG	1
1	1	1	0	0	1	1	0	Ciphertext	ľ
									Ē.

PRNG = Pseudo Random Number Generator

# Block Cyphers ("Cipher")

- Message is divided into blocks of bits
- Blocks are put through encryption functions 1 block at a time
- Suppose you are encrypting a 648-bit long message to send to your mother using a block cypher that uses 12 bits
- Your message would be chopped up into 54 blocks each 12 bits long
- Each block, in turn, would be run through a series of encryption functions (substitution and transposition)
- Ending up with 54 blocks of ciphertext



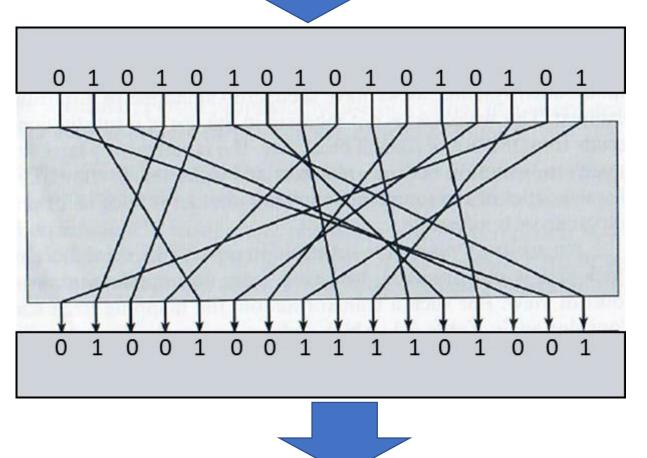
Note: Block ciphers typically use 64, 128, 512 bits at a time

Remember this example:

0101010101010101

#### 2 main attributes combined in a cypher

- 1. Confusion: usually carried out through substitution
- **2. Diffusion:** Usually carried out through <u>transposition</u>



Harris, S. and Maymi, F. (2016) <u>All-In-One CISSP Exam Guide</u>, McGraw Hill Education

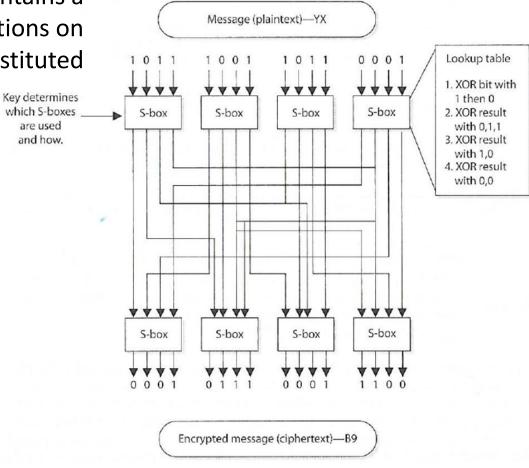
0100100111101001

### Block Ciphers versus Stream Ciphers

#### In contrast, block ciphers encrypt a block of bits at a time

In this example, each Substitution Box (S-box) contains a lookup table used by the algorithm as instructions on how the bits are substituted

Plaintext	Ciphertext	Ciphertext	Plaintext
0000	1110	0000	1110
0001	0100	0001	-0011
0010	1101	0010	0100
0011	0001	0011	1000
0100	0010	0100	0001
0101	1111	0101	1100
0110	1011	0110	1010
0111	1000	0111	1111
1000	0011	1000	0111
1001	1010	1001	1101
1010	0110	1010	1001
1011	1100	1011	0110
1100	0101	1100	1011
1101	1001	1101	0010
1110	0000	1110	0000
1111	0111	1111	0101



...followed by transposition...

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#### Encryption table

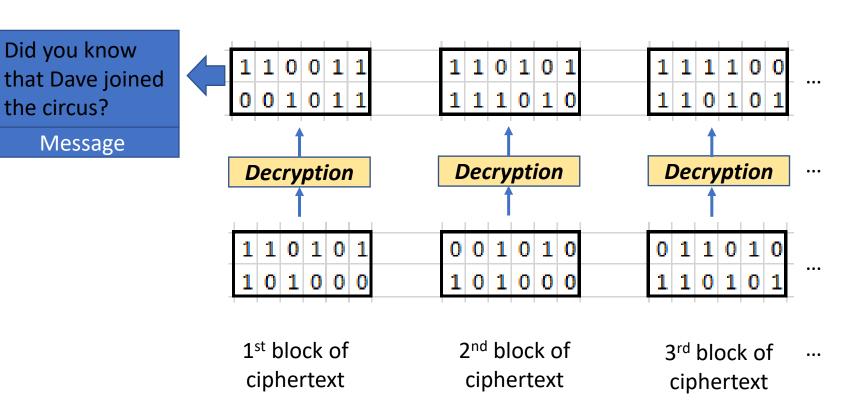
Decryption table

# Block Ciphers

- Message is divided into blocks of bits
- Blocks are put through mathematical functions 1 block at a time

You send the message to your mother. She uses the same block cipher and key (symmetric) to decipher the message

- The 54 ciphertext blocks go back through the algorithm in the reverse sequence
- Resulting in your original plaintext message your message



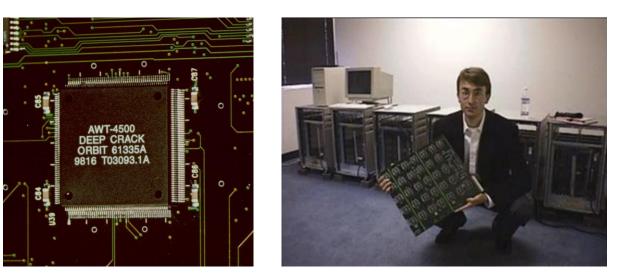
### Modern Block Ciphers

- Use block sizes of 128-bits or greater
  - Examples of Block Ciphers that can be used are:
    - <u>AES</u> (NIST's 2001 <u>Advanced Encryption Standard</u> originally known as Rijndael)
      - 128 bit block size, but 3 different key lengths: 128, 192, and 256 bits
    - Blowfish
    - Twofish
    - Serpent
- Do not use these examples of block ciphers which have a 56 bit key length, which is too small to provide secure encryption:
  - DES (Data Encryption Standard)
  - 3DES

# **Practical Cryptanalysis**

#### DES Cracker:

- A DES key search machine
- Contains 1,536 chips
- Cost: \$250,000
- Searches 88 billion keys per second
- In 1998, Electronic Freedom Foundation (EFF) won \$10,000 in RSA Laboratory's "DES Challenge II-2" by successfully finding a DES key through brute force in 56 hours



- DES uses a 56-bit key, meaning that there are 256 possible keys under which a message can be encrypted.
  - This is: 72,057,594,037,927,936, or approximately 72 quadrillion possible keys
- One of the major criticisms of DES, when proposed in 1975, was that the key size was too short
  - Martin Hellman and Whitfield Diffie of Stanford University estimated that a machine fast enough to test that many keys in a day would have cost about \$20 million in 1976,
    - An affordable sum to national intelligence agencies
- Subsequent advances in the price/performance of chips kept reducing that cost until, twenty years later, it became affordable for even a small nonprofit organization such as the EFF to mount a realistic attack

### Agenda

#### ✓ Updated Schedule

- ✓ Symmetric cryptography
- ✓ Block versus Stream ciphers
- ✓ Block ciphers
- Block ciphers mode of operations
- Hashes

### Block cipher's "mode of operation"

5 modes of operation are used to tailor them for use in different applications:

- 1. ECB Electronic Code Book mode
- 2. CBC Cipher Block Chaining mode
- 3. CFB Cipher FeedBack mode
- 4. OFB Output FeedBack mode
- 5. CTR CounTeR mode

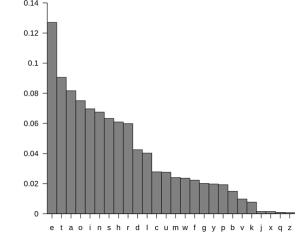
# ECB – Electronic Code Book mode

• A data block of a certain size (e.g. 64 bits or 128 bits or...) is entered into the algorithm with the key, and a block of cipher text is produced

C<sub>*i*</sub> = *Encrypt*(*Key*, *P<sub>i</sub>*) for *i* = 1, ..., *k* 

Where:

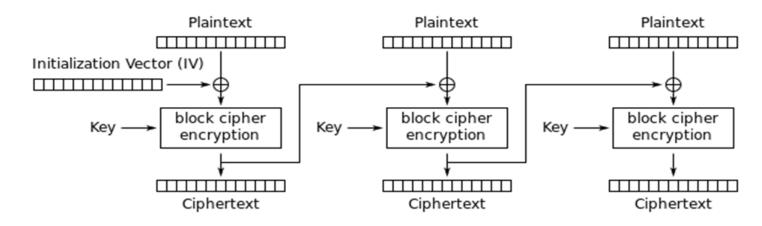
- Ci is block i of ciphertext
- P, is a block of plaintext



- Encrypts every block the same way every time for a given key
- Why is this a problem?
  - This is a problem because **frequency analysis** of the encrypted text can reveal a lot of information
  - Not enough randomness MIS 4596

# Solution: CBC – Cipher Block Chaining mode

- Is much more secure
- Does not reveal a pattern of encryption for frequency analysis
- Each block of text, the key, and the value based on the previous block are processed in the algorithm and applied to the next block of text



- XORs a plaintext with the **last** encrypted block before encrypting it. This ensures that the same plaintext is encrypted differently every time.
- Requires an initialization vector (or IV) to get started, since the first block doesn't have a previous encrypted block to XOR against.

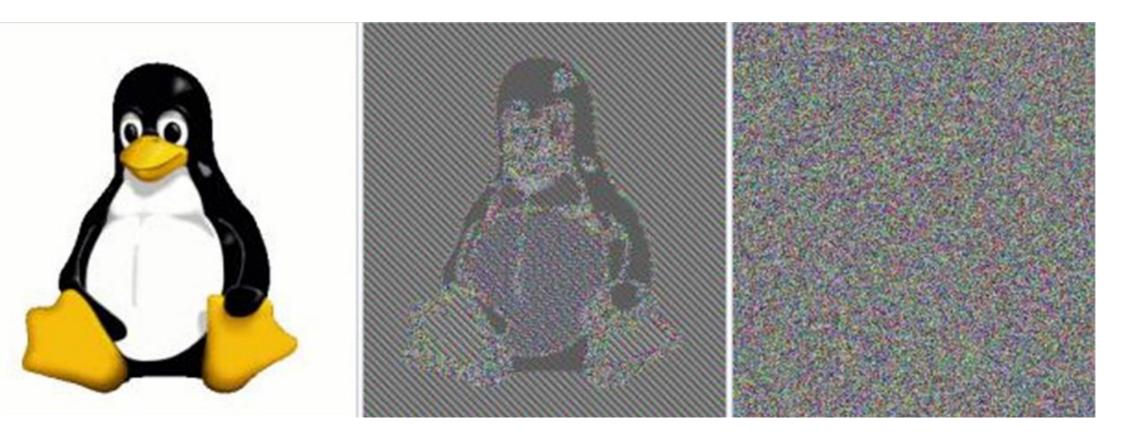
# A similar concept to *diffusion* is known as the *The Avalanche Effect*

A changed to a single plaintext bit should have an influence over several of the resulting ciphertext bits

In a strong block cipher, if 1 plaintext bit is changed, it will change every ciphertext bit with the probability of 50%

That is, if 1 plaintext bit changes, then about ½ of the ciphertext bits will change

**Avalanche Effect:** A small change to the key or to the plaintext should cause drastic changes to the resulting ciphertext



#### Original Image

Block cipher with ECB (Electronic Code Book) encryption Block cipher with CBC (Cipher Block Chaining) or any of the other modes of encryption

Not good!

These are good!

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### Agenda

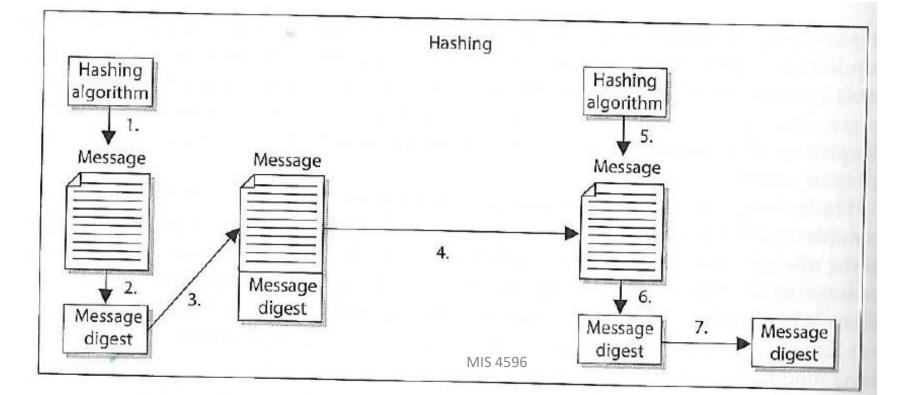
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# One-way Hash function

- Assures message integrity
- A function that takes a variable-length string (i.e. message) and produces a fixed-length value called a hash value
- Does not use keys

- 1. Sender puts message through hashing function
- 2. Message digest generated
- 3. Message digest appended to the message
- 4. Sender sends message to receiver
- 5. Receiver puts message through hashing function
- 6. Receiver generates message digest value
- Receiver compares the two message digests values.
  If they are the same, the message has not been altered

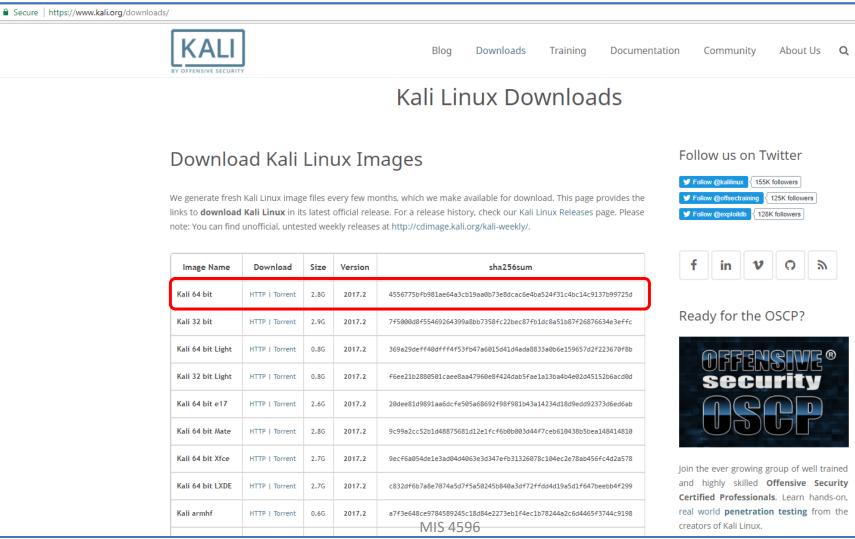


#### Hashing results in **fixed-sized output**

Names for the output of a hashing functions include "hash" and a *message digest* (md), because a hash "digests" an input of any size down to a **fixed-sized output** 

- No matter the size of the input, the output is the same, for example the md5 hash function's output:
  - Letter 'a' in binary: 01000001 => md5 hash => 32-character string
  - Blu-ray disk digest => md5 hash => 32-character string
  - 6 TB hard drive digest => md5 hash => 32-character string

One useful example: Testing the integrity of a file (e.g. program) downloaded from the internet...



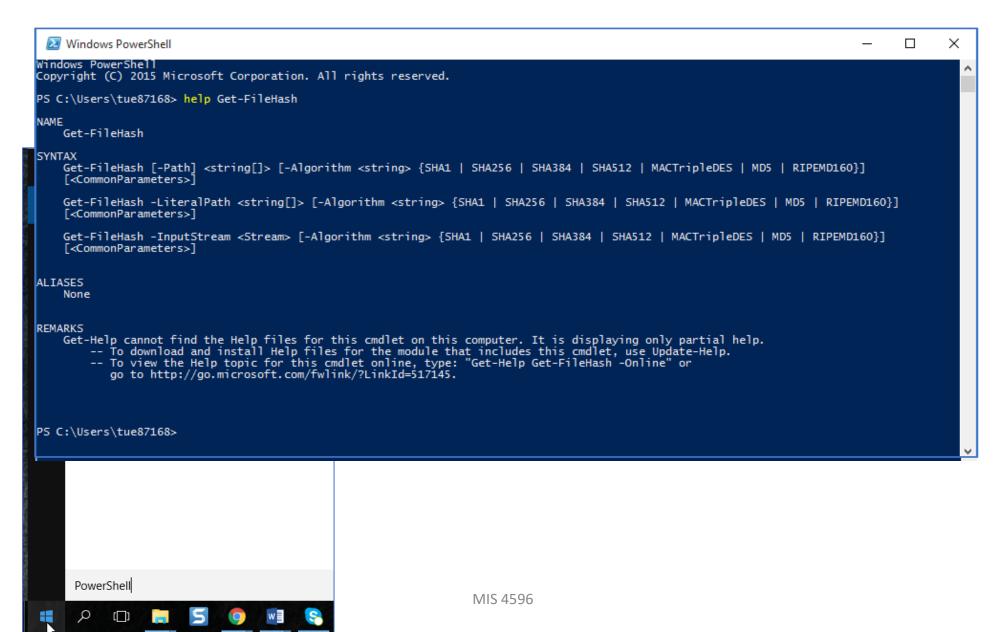
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#### *Testing the integrity of a file (e.g. program) from the internet...*

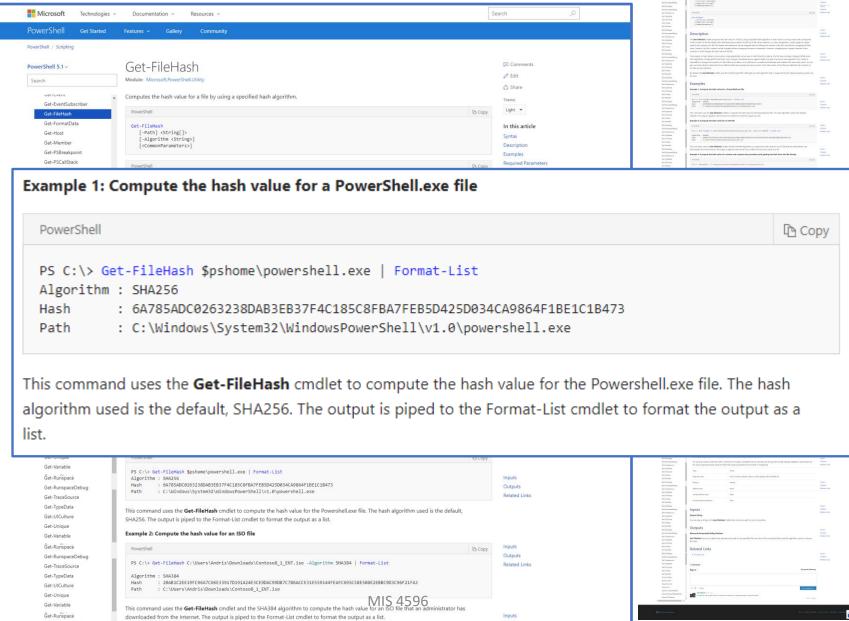
Image Name	Download	Size	Version	sha256sum
Kali 64 bit	HTTP   Torrent	2.8G	2017.2	4556775bfb981ae64a3cb19aa0b73e8dcac6e4ba524f31c4bc14c9137b99725d

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🕹 Downloads	*	170919_ITACS_Fall17_Brochure_Proof3-WithComments (1)	11/8/2017 5:49 PM	Adobe Acrobat Docu	215 KB	•
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	~	🗾 170919_ITACS_Fall17_Brochure_Proof3	11/8/2017 5:32 PM	Adobe Acrobat Docu	208 KB	
This PC	× -	Secure Components Risk and Controls Library	11/8/2017 11:34 AM	Microsoft Excel Work	103 KB	
1,370 items 1 item selected	5	1 15				

#### Is the Kali I downloaded the same Kali that was published?



https://docs.microsoft.com/en-us/powershell/module/microsoft.powershell.utility/get-filehash?view=powershell-5.1



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Windows PowerShell

Windows PowerShell Copyright (C) 2015 Microsoft Corporation. All rights reserved.

PS C:\Users\tue87168> dir

Directory: C:\Users\tue87168

Mode	LastV	VriteTime	Length	Name
 d d-r d-r d-r d-r d-r d-r d-r	Lastv  9/27/2016 8/21/2016 10/13/2017 11/5/2017 11/7/2017 10/13/2017 11/6/2017 11/6/2017 11/7/2017 11/2/2017 6/20/2017 8/10/2016	11:28 AM 10:57 AM 8:35 AM 8:48 PM 2:31 PM 8:35 AM 9:33 AM 2:53 PM 8:16 AM 5:07 PM	Length 	Name  .oracle_jre_usage Benefits Contacts Desktop Documents Downloads Favorites Google Drive Intel Links logs MIS
d-r d-r d-r d	10/13/2017 11/2/2017 11/9/2017 8/8/2016 10/13/2017 10/13/2017 11/17/2016 10/13/2017	8:35 AM 8:16 AM 11:46 AM 11:20 AM 8:35 AM 8:35 AM 11:20 AM		Music OneDrive Pictures Roaming Saved Games Searches Tracing Videos

PS C:\Users\tue87168> <mark>cd</mark> Downloads PS C:\Users\tue87168\Downloads> <mark>di</mark>r \*.iso

Directory: C:\Users\tue87168\Downloads

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-a	8/10/2017 10:	55 AM	674803712	CSET_8.0 (1).iso
-a	8/10/2017 11:	03 AM	674803712	CSET_8.0 (2).iso
-a	6/12/2017 10:	29 AM		CSET_8.0.iso
-a	9/27/2017 3:	03 PM	2421987328	en_project_professional_2016_x86_x64_dvd_6962236.iso
-a	10/3/2017 8:	49 PM	2421987328	en_visio_professional_2016_x86_x64_dvd_6962139.iso
-a	11/11/2016 11:	45 AM		Fedora-Live-Workstation-x86_64-23-10.iso
-a	11/9/2017 2:	31 PM	3020619776	kali-linux-2017.2-amd64.iso

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Image Name	Download	Size	Version	sha256sum
Kali 64 bit	HTTP   Torrent	2.8G	2017.2	4556775bfb981ae64a3cb19aa0b73e8dcac6e4ba524f31c4bc14c9137b99725d

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PS C:∖Use	rs\tue87168> <mark>cd</mark> rs\tue87168\Dow	nloads> dir				^
Mode	tory: C:\Users\ Last	UriteTime	n roads Length M	Name		
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Algorithm Hash Path	: 45567758FB98			4BA524F31C4BC14C9137B99725D x-2017.2-amd64.iso		
PS C:\Use	rs\tue87168\Dow	nloads> _		MIS 4596		~

🛃 Windo	ws PowerShell				-	- 🗆	$\times$
PS C:\Use	rs\tue87168\Down]	oads> dir *.	.txt	Notice the amount of	<u>confusion</u> and <u>d</u>	iffusion	^
Direct	tory: C:\Users\tu	e87168\Down	loads	resulting from a 1-cha	acter change!		
Mode	LastWr	iteTime	Length M	lame			
 -a	11/9/2017	3:04 PM	15 M	IIS5206-IsGood.txt			
Algorithm Hash Path				D9D8C9BD4B2A45A4ACCB3999DD Good.txt			
MIS5206 i				od.txt 206-IsGood.txt   Format-List			
Algorithm Hash Path	: SHA256 : 877B45EA5D40D9 : C:\Users\tue87	8FF8D1ABD919 168\Download	9E154F446FEA1 ds\MIS5206-Is	1387DBB13DDEE448F9932928A5 🤞			
PS C:\Use	rs\tue87168\Down]	oads>		MIS 1596			J

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# File Integrity Monitoring

- An internal control process
- Validates the integrity of operating system and application software files
- Uses hash verification to compare the current file state and a known good baseline state
- Involves calculating and storing a hash value of a known good version of the file ("original baseline")
- Compares the baseline with the calculated hash of the current state of the file to detect unauthorized changes

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EMAILS

REAL-TIME

Behaviori

FIM MANAGEMENT

SERVER

CHANGE LOGS

SERVER

STACK

#### How are passwords stored in Linux?

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#### Passwords are hashed! Don't store plaintext passwords

cat /etc/passwd root:x:0:0:root:/root:/bin/bash daemon:x:1:1:daemon:/usr/sbin:/bin/sh bin:x:2:2:bin:/bin:/bin/sh sys:x:3:3:sys:/dev:/bin/sh sync:x:4:65534:sync:/bin:/bin/sync games:x:5:60:games:/usr/games:/bin/sh man:x:6:12:man:/var/cache/man:/bin/sh lp:x:7:7:lp:/var/spool/lpd:/bin/sh mail:x:8:8:mail:/var/mail:/bin/sh news:x:9:9:news:/var/spool/news:/bin/sh uucp:x:10:10:uucp:/var/spool/uucp:/bin/sh proxy:x:13:13:proxy:/bin:/bin/sh www-data:x:33:33:www-data:/var/www:/bin/sh backup:x:34:34:backup:/var/backups:/bin/sh list:x:38:38:Mailing List Manager:/var/list:/bin/sh irc:x:39:39:ircd:/var/run/ircd:/bin/sh gnats:x:41:41:Gnats Bug-Reporting System (admin):/var/lib/gnats:/bin/sh nobody:x:65534:65534:nobody:/nonexistent:/bin/sh libuuid:x:100:101::/var/lib/libuuid:/bin/sh dhcp:x:101:102::/nonexistent:/bin/false syslog:x:102:103::/home/syslog:/bin/false klog:x:103:104::/home/klog:/bin/false sshd:x:104:65534::/var/run/sshd:/usr/sbin/nologin msfadmin:x:1000:1000:msfadmin,,,:/home/msfadmin:/bin/bash bind:x:105:113::/var/cache/bind:/bin/false postfix:x:106:115::/var/spool/postfix:/bin/false ftp:x:107:65534::/home/ftp:/bin/false postgres:x:108:117:PostgreSQL administrator,,,:/var/lib/postgresql:/bin/bash mysql:x:109:118:MySQL Server,,,:/var/lib/mysql:/bin/false tomcat55:x:110:65534::/usr/share/tomcat5.5:/bin/false distccd:x:111:65534::/:/bin/false user:x:1001:1001:just a user,111,,:/home/user:/bin/bash service:x:1002:1002:,,,:/home/service:/bin/bash telnetd:x:112:120::/nonexistent:/bin/false proftpd:x:113:65534::/var/run/proftpd:/bin/false statd:x:114:65534::/var/lib/nfs:/bin/false

cat /etc/shadow root:\$1\$/avpfBJ1\$x0z8w5UF9Iv./DR9E9Lid.:14747:0:99999:7::: daemon:\*:14684:0:99999:7::: bin:\*:14684:0:99999:7::: sys:\$1\$fUX6BPOt\$Miyc3UpOzQJqz4s5wFD910:14742:0:99999:7::: sync:\*:14684:0:99999:7::: games:\*:14684:0:99999:7::: man:\*:14684:0:99999:7::: lp:\*:14684:0:99999:7::: mail:\*:14684:0:99999:7::: news:\*:14684:0:99999:7::: uucp:\*:14684:0:99999:7::: proxy:\*:14684:0:99999:7::: www-data:\*:14684:0:99999:7::: backup:\*:14684:0:99999:7::: list:\*:14684:0:99999:7::: irc:\*:14684:0:99999:7::: gnats:\*:14684:0:99999:7::: nobody:\*:14684:0:99999:7::: libuuid:!:14684:0:99999:7::: dhcp:\*:14684:0:99999:7::: syslog:\*:14684:0:99999:7::: klog:\$1\$f2ZVMS4K\$R9XkI.CmLdHhdUE3X9jqP0:14742:0:99999:7::: sshd:\*:14684:0:99999:7::: msfadmin:\$1\$XN10Zj2c\$Rt/zzCW3mLtUWA.ihZjA5/:14684:0:99999:7::: bind:\*:14685:0:99999:7::: postfix:\*:14685:0:99999:7::: ftp:\*:14685:0:99999:7::: postgres:\$1\$Rw35ik.x\$MgQgZUuO5pAoUvfJhfcYe/:14685:0:99999:7::: mysql:!:14685:0:99999:7::: tomcat55:\*:14691:0:99999:7::: distccd:\*:14698:0:99999:7::: user:\$1\$HESu9xrH\$k.o3G93DGoXIiQKkPmUgZ0:14699:0:99999:7::: service:\$1\$kR3ue7JZ\$7GxELDupr50hp6cjZ3Bu//:14715:0:99999:7::: telnetd:\*:14715:0:99999:7::: proftpd:!:14727:0:99999:7::: statd:\*:15474:0:99999:7:::

#### Shadow file

# Cryptanalysis Attacks

- Brute force
  - Trying all key values in the keyspace
- Frequency Analysis
  - Guess values based on frequency of occurrence
- Dictionary Attack
  - Find plaintext based on common words
- Known Plaintext
  - Format or content of plaintext available
- Chosen Plaintext
  - Attack can encrypt chosen plaintext
- Chosen Ciphertext
  - Decrypt known ciphertext to discover key

- Random Number Generator (RNG) Attack
  - Predict initialization vector used by an algorithm
- Social Engineering
  - Humans are the weakest link

# Cryptanalysis Attacks

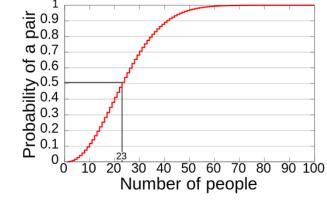
Collisions

- Two different messages with the same hash value
- Based on the "birthday paradox"
- Hash algorithms should be resistant to this attack

The birthday paradox, also known as the birthday problem, states that in a random group of 23 people, there is about a 50 percent chance that two people have the same birthday.

# Is the Birthday Attack Real?

There are multiple reasons why this seems like a paradox



One is that when in a room with 22 other people, if a person compares her/his birthday with the birthdays of the other people it would make for only 22 comparisons—only 22 chances for people to share the same birthday.

When all 23 birthdays are compared against each other, it makes for much more than 22 comparisons. How much more?

Well, the first person has 22 comparisons to make, but the second person was already compared to the first person, so there are only 21 comparisons to make.

The third person then has 20 comparisons, the fourth person has 19 and so on.

If you add up all possible comparisons (22 + 21 + 20 + 19 + ... +1) the sum is 253 comparisons, or combinations. Consequently, each group of 23 people involves 253 comparisons, or 253 chances for matching birthdays.

### MD5 (Message Digest 5)

- A 128-bit hash algorithm, still in common use
- Has been broken
- 128-bit hash, but only need  $2^{128/2} = 2^{64}$  to find a collision
- Not strong enough for modern computers

### SHA -1 (Security Hash Algorithm 1)

- A 160-bit hash algorithm, still in common use
- Has been broken
- 160-bit hash, but only need  $2^{160/2} = 2^{80}$  to find a collision
- No longer strong enough for modern computers

🔀 Windows Po	owerShell	- 🗆 X
PS C:\Users\D	ave\Desktop\MD5-Hash-Collision-Example> get-filehash ProgramA.exe -Alg	orithm SHA256
Algorithm	Hash	Path 
SHA256	60D13913155644883F130B85EB24D778314014C9479AEDB5F6323BF38AD3A451	C:\Users\Dave\Desktop\MD5-Hash-Collision-Example\ProgramA.exe
PS C:\Users\D	ave\Desktop\MD5-Hash-Collision-Example> get-filehash ProgramB.exe -Alg	orithm SHA256
Algorithm	Hash	Path 
SHA256	1316543942A8C6CD754855500CD37068EDBBD8B31C4979D2825A4E799FED6102	C:\Users\Dave\Desktop\MD5-Hash-Collision-Example\ProgramB.exe

#### Hello, world!

This program is evil!!! Erasing hard drive...1Gb...2Gb... just kidding! Nothing was erased.

ProgramA run

(press enter to quit)

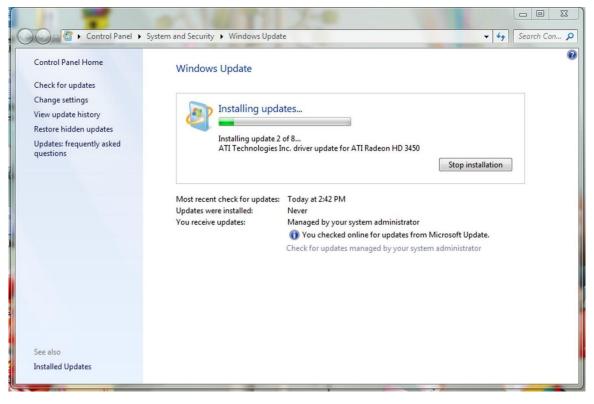
(press enter to quit)

#### ProgramB run

🔀 Windows P	owerShell	- 0
PS C:\Users\[	<pre>Dave\Desktop\MD5-Hash-Collision-Example&gt; get-filehash Pr</pre>	rogramA.exe -Algorithm MD5
Algorithm	Hash	Path
MD5	CDC47D670159EEF60916CA03A9D4A007	C:\Users\Dave\Desktop\MD5-Hash-Collision-Example\ProgramA.exe
PS C:\Users\[	Dave\Desktop\MD5-Hash-Collision-Example> get-filehash Pr	rogramB.exe -Algorithm MD5
Algorithm	Hash	Path

The malware Flame used a MD5 hash collision to hijack Microsoft Windows Update and spread itself across networks

- Flame collected audio, keystrokes, screenshots which it sent to a malicious server
- Found a collision within a single millisecond
- Cost ~\$200k computing time just for 1ms
- Attributed to advanced persistent threat group Equation Group
- Espionage attacks on countries in and around Iran



#### PWN (verb)

1. An act of dominating an opponent.

2. Great, ingenious; applied to methods and objects.

- Originally dates back to the days of WarCraft, when a map designer mispelled "Own" as "Pwn"
- What was originally suppose to be "player has been owned." was "player has been pwned"

Use of the term "Pwn' grew and is now used throughout the online world, especially in online games:

- 1. "I pwn these guys on battlenet"
- 2. "This strategy pwns!" or "This game pwn."

https://www.urbandictionary.com/define.php?term=pwn



Equation Group's Flame malware won 2012 "Epic Ownage" Pwnie award

- Pwnie Awards recognize both excellence and incompetence in the field of information security
- Awards are presented yearly at the Black Hat Security Conference

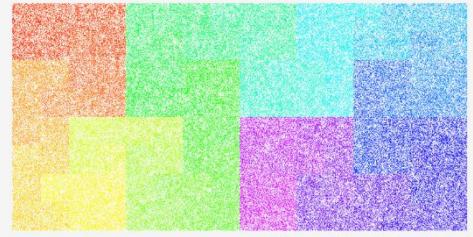
Hashing algorithms are used for browser ssl (secure sockets layer)

- In 2014, many sites were still using SHA-1, at the time known to be dangerously vulnerable
- Google declared state of emergency to push companies to upgrade

ABOUT BLOG RESUME

#### Why Google is Hurrying the Web to Kill SHA-1

published by Eric Mill on September 7, 2014, 58 comments



Hilbert map of hashing algorithms, by lan Boyd

Most of the secure web is using an insecure algorithm, and Google's just declared it to be a slow-motion emergency.

Something like <u>90% of websites</u> that use SSL encryption — https:// — use an algorithm called <u>SHA-1</u> to protect themselves from being impersonated. This guarantees that when you go to https://www.facebook.com, you're visiting the real Facebook and not giving your password to an attacker.

Unfortunately, <u>SHA-1 is dangerously weak</u>, and has been for a <u>long time</u>. It gets weaker every year, but remains widely used on the internet. Its replacement, <u>SHA-2</u>, is strong and supported just about everywhere.

Google <u>recently announced</u> that if you use Chrome, then you're about to start seeing a progression of warnings for many secure websites:



What's <u>about to befall websites</u> with SHA-1 certificates that expire in 2017, in Chrome.

SHA-2 uses 224, **256**, 384, and 512-bit hashes

- But... it is built using the design of SHA-1, and prone to the same weaknesses
- It's believed to be a matter of time before SHA-2 is also exploited
- SHA-3 was just ratified recently by NIST, the U.S. National Institute of Standards and Technology
  - It was the result of a six-year hashing competition. Also uses 224-, 256-, 384-, 512-bit hashes

#### Why does this matter for businesses?

Business needs a reliable way to prove integrity of data, files, programs, that can be trusted

### Agenda

#### ✓ Updated Schedule

- ✓ Symmetric cryptography
- ✓ Block versus Stream ciphers
- ✓ Block ciphers
- $\checkmark$  Block ciphers mode of operations
- ✓ Hashes